

Lightweight Memory Tracing

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Memory Tracing via Memlets

Execute code (*memlets*) for every memory access

A memlet inspects a single memory access based on target *address*, *type* of memory access, *instruction*, or prior *state*

Memory tracing enables detailed memory access logs, debugging of memory accesses, security checks, privacy extensions

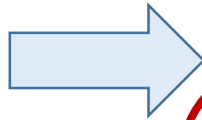
Memory Tracing by Example

Binary translation weaves memlets into executed code

memTrace is general, for talk let's focus on example:

- Unlimited *watchpoints*: check if R/W watchpoint is set

```
addl (%ebx), %eax  
jg bb1  
jmp bb2
```



```
/* check */  
lea (%ebx), %reg  
cmpl 0xshadow(%reg), $0x0  
jnz handler_92746  
/* translated instruction */  
addl (%ebx), %eax  
jg bb1  
jmp bb2
```

Key to Lightweight Memory Tracing

Modern CPUs support multiple ISAs: x86/x86_64

- Most programs still 32-bit x86

Cross-ISA binary translation allows the tracer to use additional hardware available in target ISA:

- Wider address space: isolation & performance
- Additional registers: flexibility & performance

Outline

Motivation and Introduction

Lightweight Memory Tracing

- Requirements
- User-defined Memlets
- Cross-ISA Binary Translation (BT)
- Implementation

Evaluation

Related Work

Conclusion

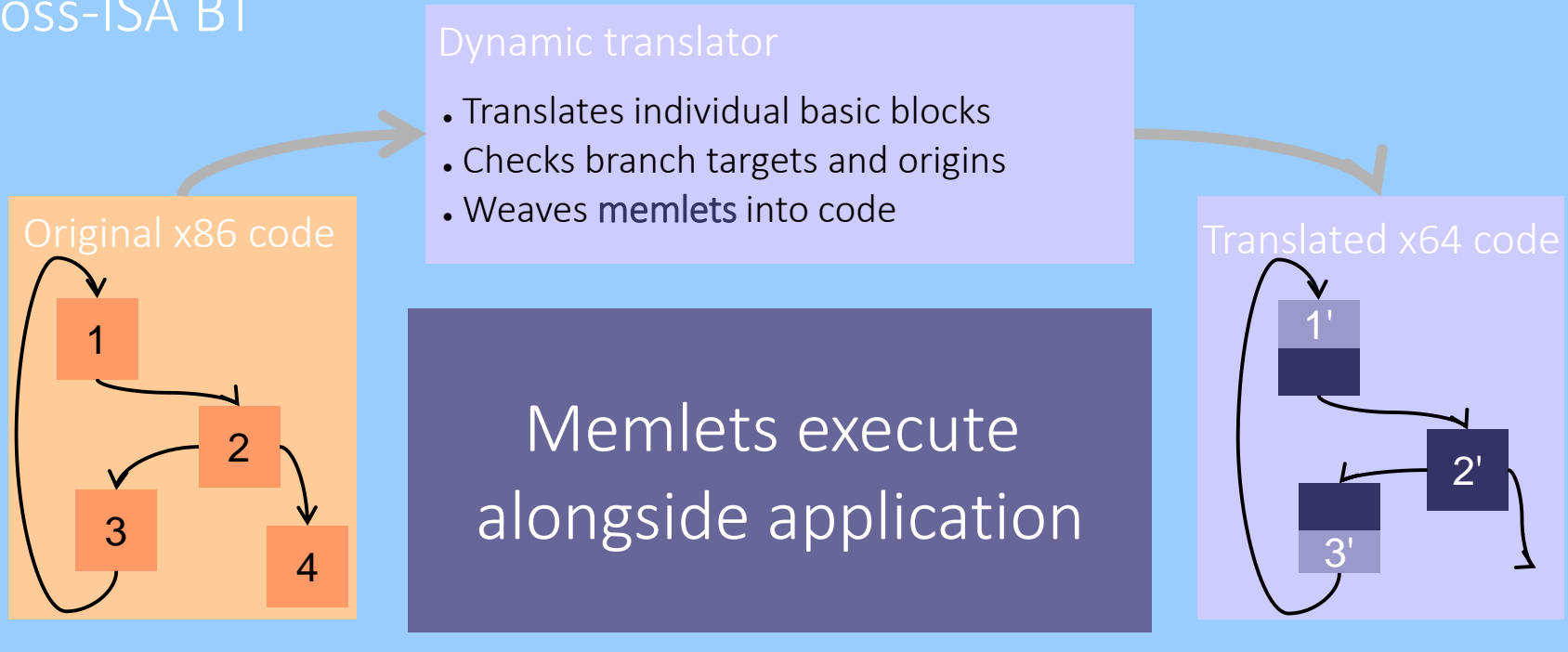
Tracing Requirements



Flexibility through BT

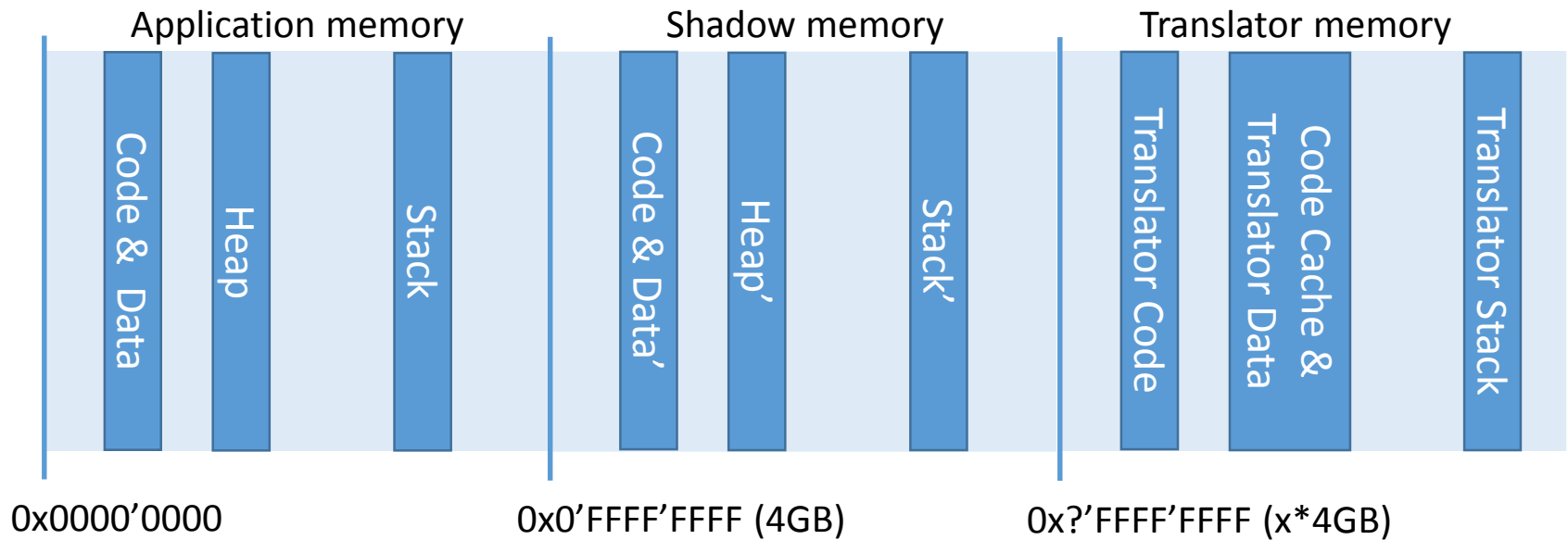


Cross-ISA BT



x64 Kernel

Isolation: Larger Memory Space



Wider memory space
Isolates tracer from application

Key to Low Overhead



Fast, efficient binary translation

Letting the hardware do most of the work...

- use 64-bit addressing (aligned 4GB blocks)
- keep state in additional/wider registers
- optimize for EFLAGS usage

Implementation



memTrace implementation (open source)

- Cross-ISA translator
- Sample memlets

Small, lean implementation

	Code	Comments
memTrace	13,800*	3,300
Memlets	150-200	100-200

*4,900 LOC for the translation tables

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Lightweight Memory Tracing

Evaluation

- Unlimited Watchpoints
- Safe Memory Allocation

Related Work

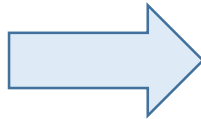
Conclusion

Unlimited Watchpoints

Watchpoints trigger on memory reads/writes

Memlet checks if read/write watchpoint is set for each memory access

```
addl (%ebx), %eax  
jg bb1  
jmp bb2
```



```
/* check */  
lea (%ebx), %r8  
cmpl 0x100000000(%r8), $0x0  
jnz handler_92746  
/* translated instruction */  
addl (%ebx), %eax  
jg bb1  
jmp bb2
```

Evaluation Setup

SPEC CPU2006 benchmarks evaluated

- System: Ubuntu 12.04, GCC 4.6.3 (64bit)
- Intel Core i7-2640M @ 2.80GHz, 4GB RAM

Four configurations:

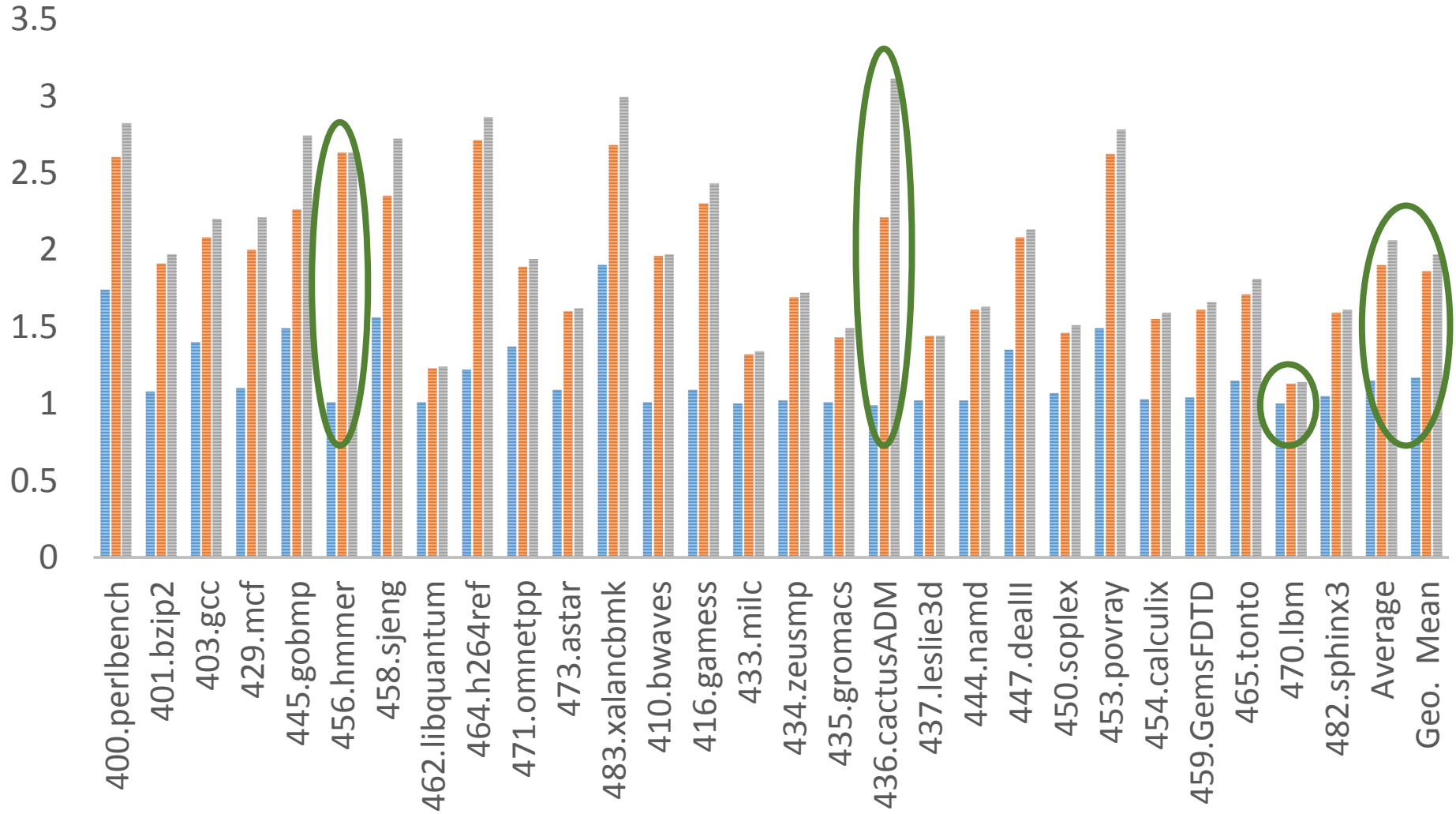
- Native
- Binary translation (BT) only
- Memory Tracing
- Full Watchpoints

SPEC CPU 2006: Low Perf. Impact

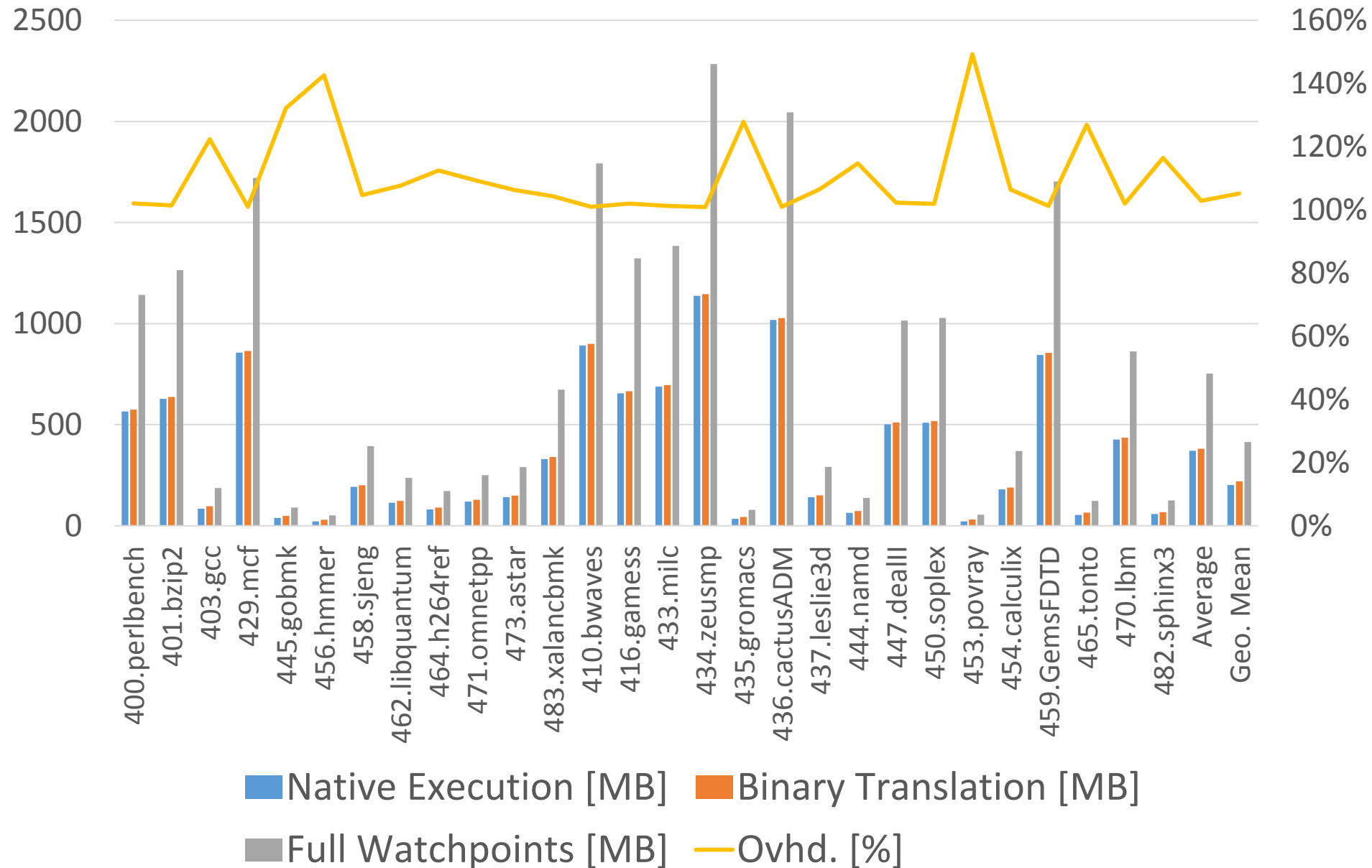
Binary Translation

Memory Tracing

Full Watchpoints



Memory Overhead: 2x



Safe Memory Allocation

Check for use-after-free bugs and heap corruption

Intercept calls to `malloc` and `free`

- Protect metadata of allocated blocks
- Check for read/write accesses to freed blocks until they are reused

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Related work

Valgrind allows high-level transformations on machine code with performance cost ($\sim 7x$ for nullgrind, $\sim 26x$ for memcheck)

GDB/Hardware watchpoints allow a limited set of watchpoints with negligible overhead

Limitations of other dynamic tracing systems are (i) limited ISA support, (ii) high overhead, or (iii) limited flexibility

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memTrace enables lightweight, low-overhead <90% memory inspection for unmodified applications

- Use resources of modern CPUs

Memlets allow user-configurable checks for each memory access

- Flexible framework for memory tracing

Source:

- <http://nebelwelt.net/projects/memTrace/>
- <https://github.com/gannimo/memTrace>