

String Oriented Programming

Exploring Format String Attacks

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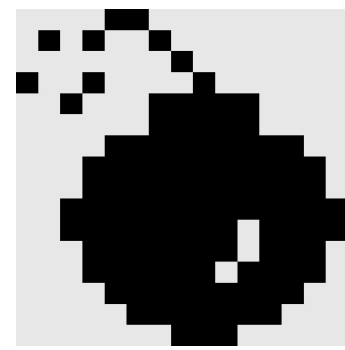


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Motivation

- Additional protection mechanisms prevent many existing attack vectors
- Format string exploits are often overlooked
 - Drawback: hard to construct (new protection mechanisms)
 - Define a way to deterministically exploit format string bugs

Attack model



- Attacker with restricted privileges forces escalation
- Attacker knows source code and binary
- Successful attacks
 - Redirect control flow to alternate location
 - Injected code is executed or alternate data is used for existing code

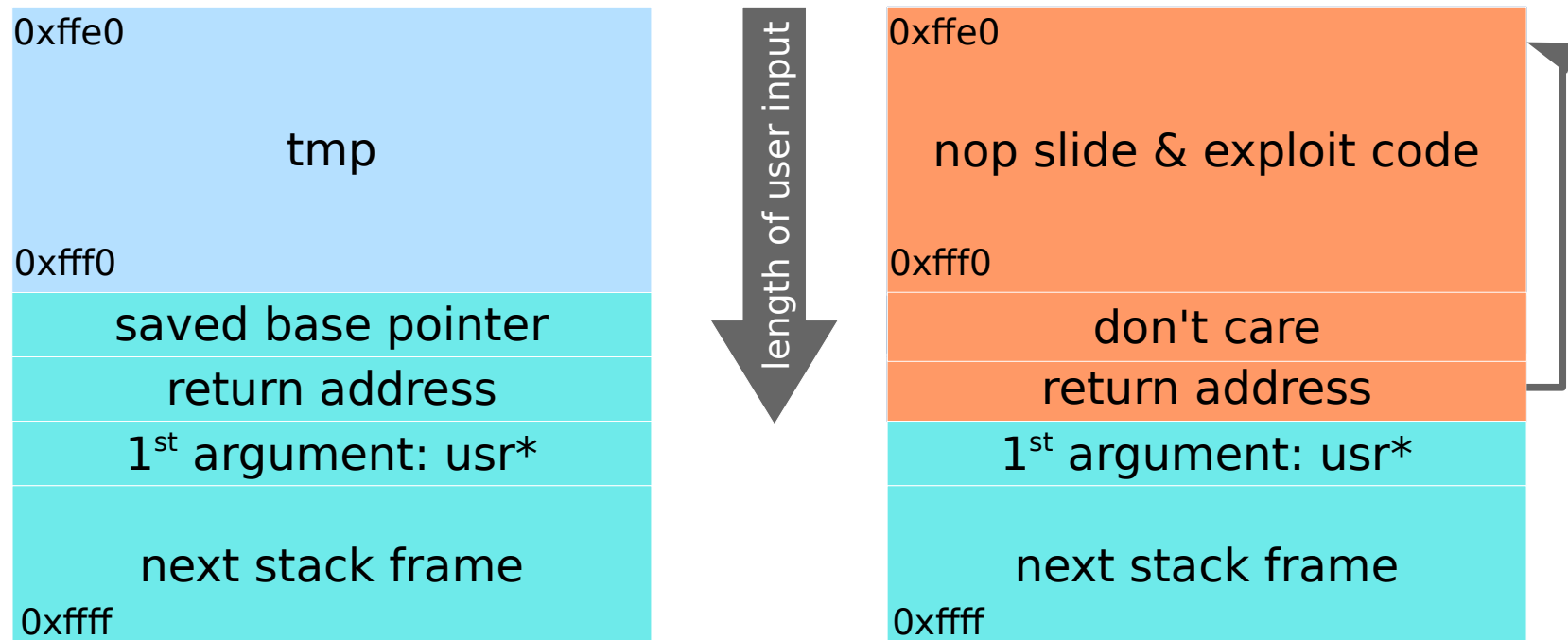
Outline

- Motivation
- Attack model
- Attack vectors and protection mechanisms
- String Oriented Programming
- Conclusion

Code injection*

- Injects additional code into the runtime image
 - Buffer overflow used to inject code as data

```
void foo(char *usr)
{
    char tmp[len];
    strcpy(tmp, usr);
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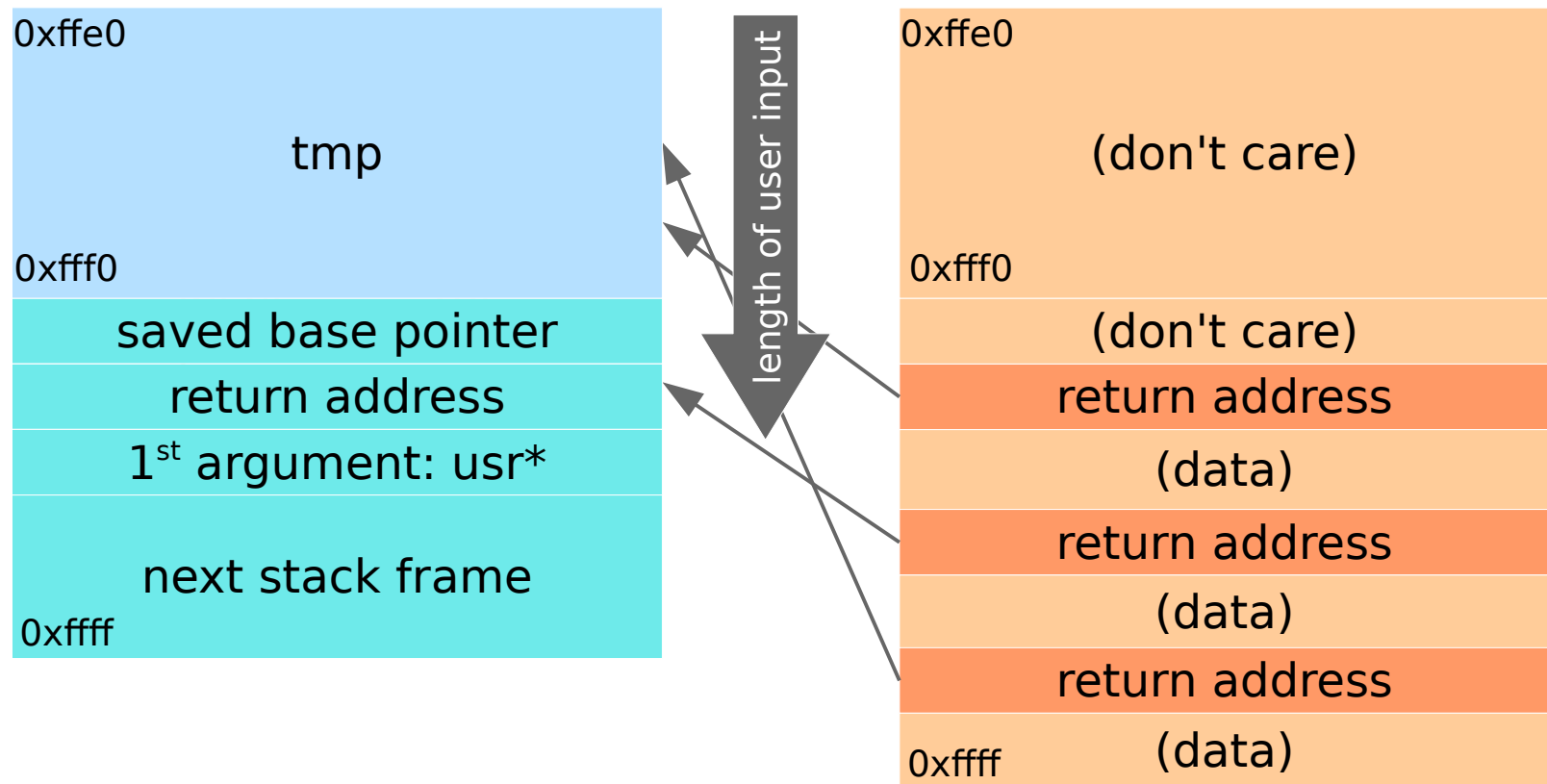
- Modern hardware and operating systems separate data and code
 - Code injection is no longer feasible due to $W \oplus X$
 - If the attacked program uses a JIT then WX pages might be available

Protection mechanisms

- Data Execution Prevention (DEP / ExecShield)
 - Enforces the executable bit ($W \oplus X$) on page granularity
 - Changes: HW, kernel, loader
- Address Space Layout Randomization (ASLR)
 - All memory addresses (heap / stack / libraries) are dynamic
 - Application itself is static
 - Changes: loader
- ProPolice (in gcc)
 - Uses canaries on the stack to protect from stack-based overflows
 - Changes: compiler

Return Oriented Programming (ROP)*

- ROP prepares several stack invocation frames
 - Executes arbitrary code
 - Stack-based buffer overflow as initial attack vector

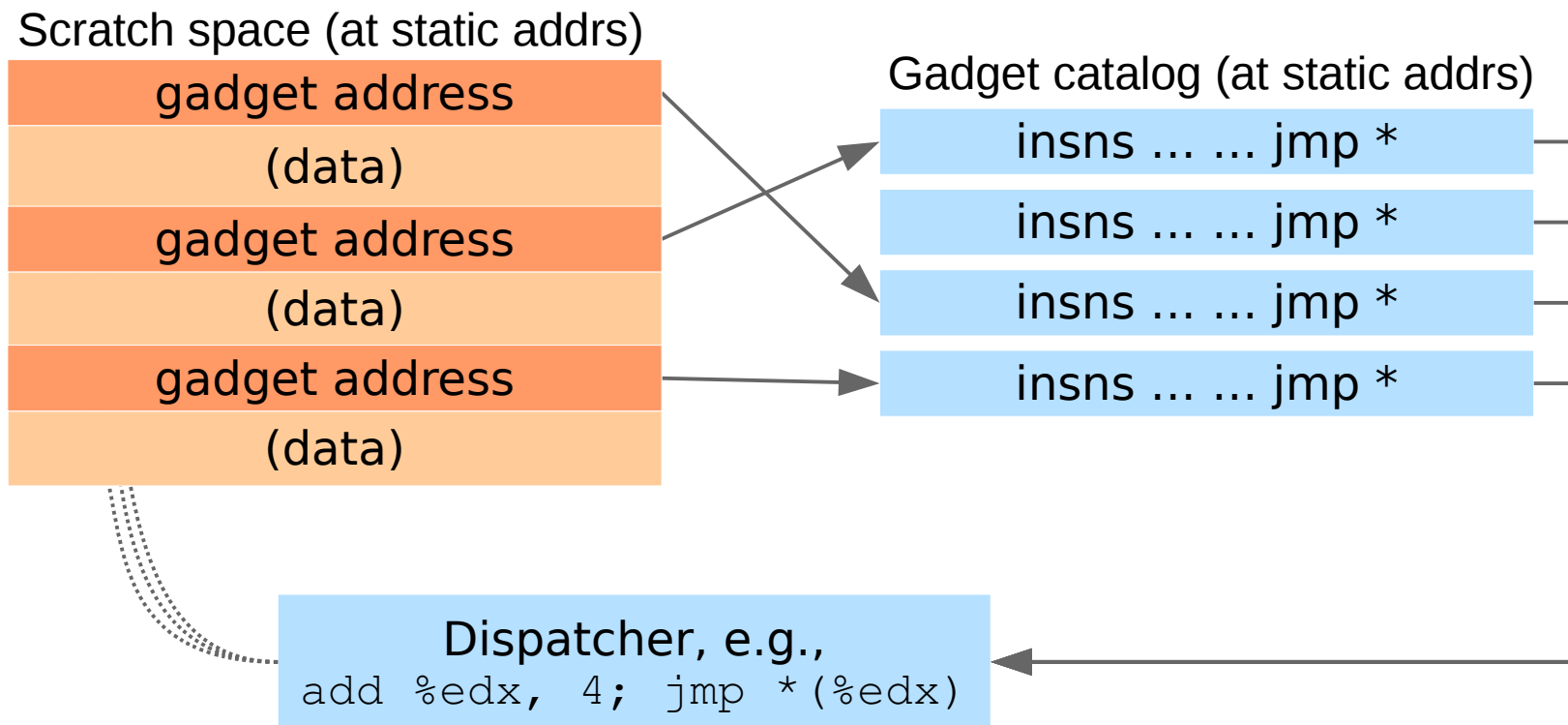


Return Oriented Programming (ROP)*

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- Executes alternate data with existing code
 - Circumvents $W \oplus X$
 - Hard to get around ASLR, ProPolice

Jump Oriented Programming (JOP)*

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 - JOP extends and generalizes ROP
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 - Hard to get around ASLR, ProPolice (if stack data used)

Format string attack*

- Attacker controlled format results in random writes
 - Format strings consume parameters on the stack
 - %n token inverses order of input, results in indirect memory write
 - Often string is on stack and can be used to store pointers

```
printf(  
    "AAAACAAA"           /* encode 2 halfword pointers */  
    "%1$49387c"         /* write 0xc0f3 - 8 bytes */  
    "%6$hn"             /* store at second HW */  
    "%1$61204c%5$hn"    /* repeat with 0xb007 */  
);
```

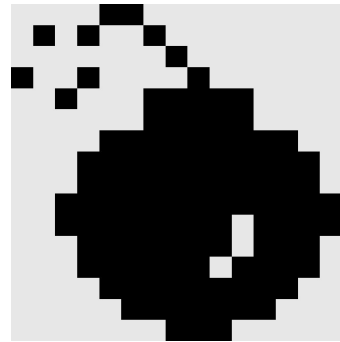
- Random writes are used to:
 - Redirect control flow
 - Prepare/inject malicious data

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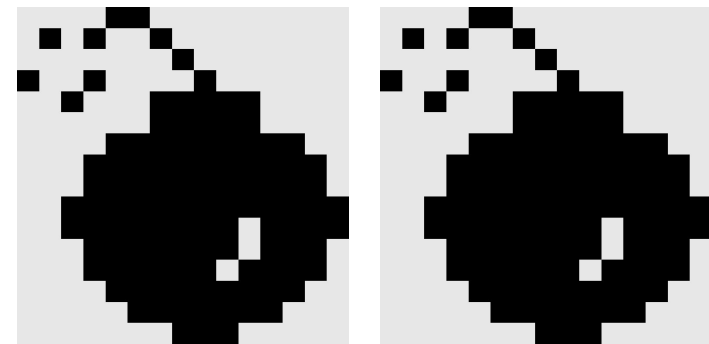
String Oriented Programming (SOP)

- SOP executes arbitrary code (through data)
 - Needed: format string bug, attacker-controlled buffer on stack
 - Not needed: buffer overflow, executable memory regions
- Executing code
 - SOP builds on ROP/JOP
 - Overwrites static instruction pointers (to initial ROP/JOP gadgets)



String Oriented Programming

- SOP patches and resolves addresses
 - Application is static (this includes .plt and .got)
 - Static program locations used to resolve relative addresses
- Resolving hidden functions
 - ASLR randomizes ~10bit for libraries
 - Modify parts of static .got pointers
 - Hidden functions can be called without loader support

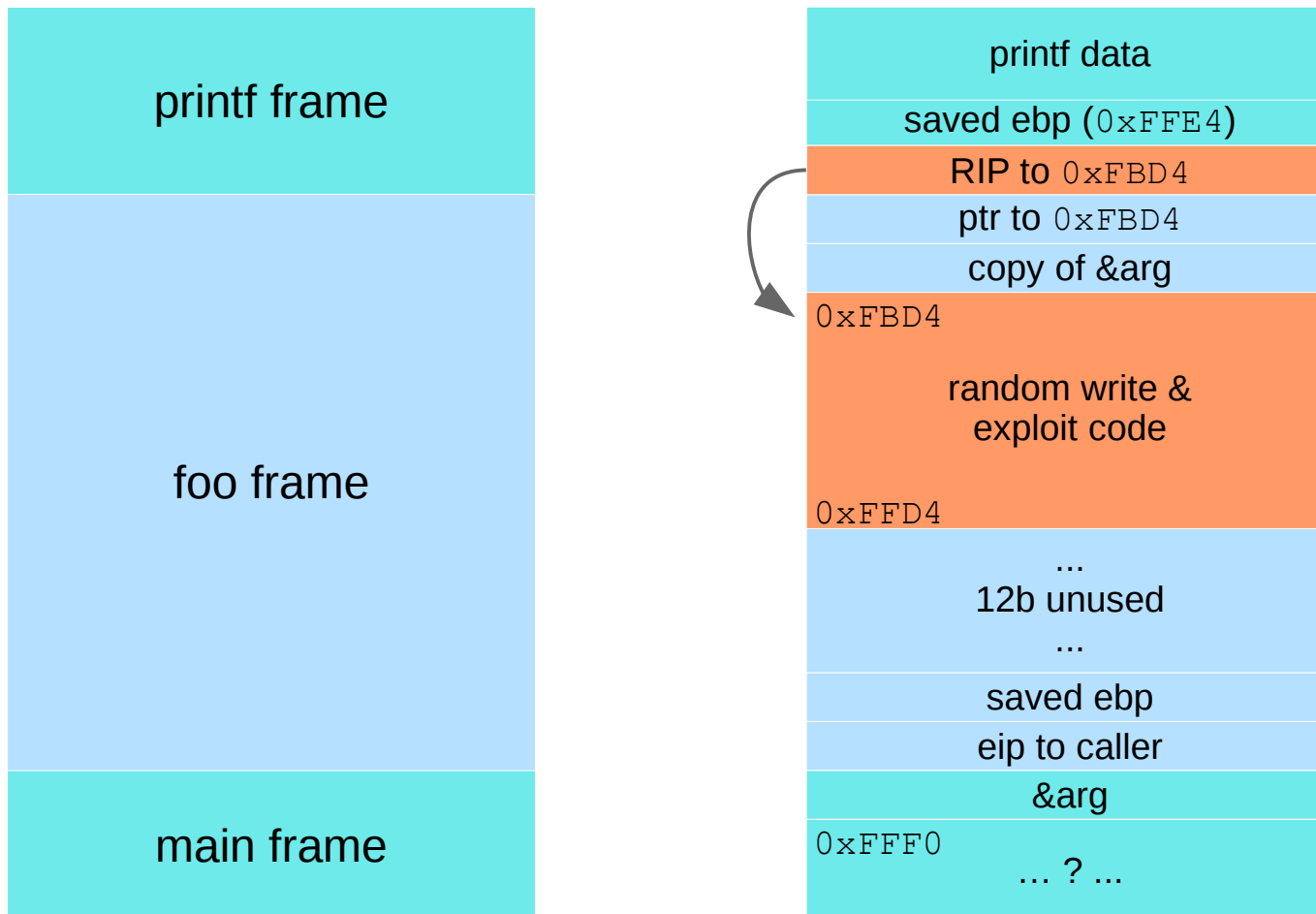


Running example

```
void foo(char *arg) {  
    char text[1024];           // buffer on stack  
    if (strlen(arg) >= 1024) // length check  
        return;  
    strcpy(text, arg);  
    printf(text);             // vulnerable printf  
}  
  
...  
  
foo(user_str);               // unchecked user data  
  
...
```

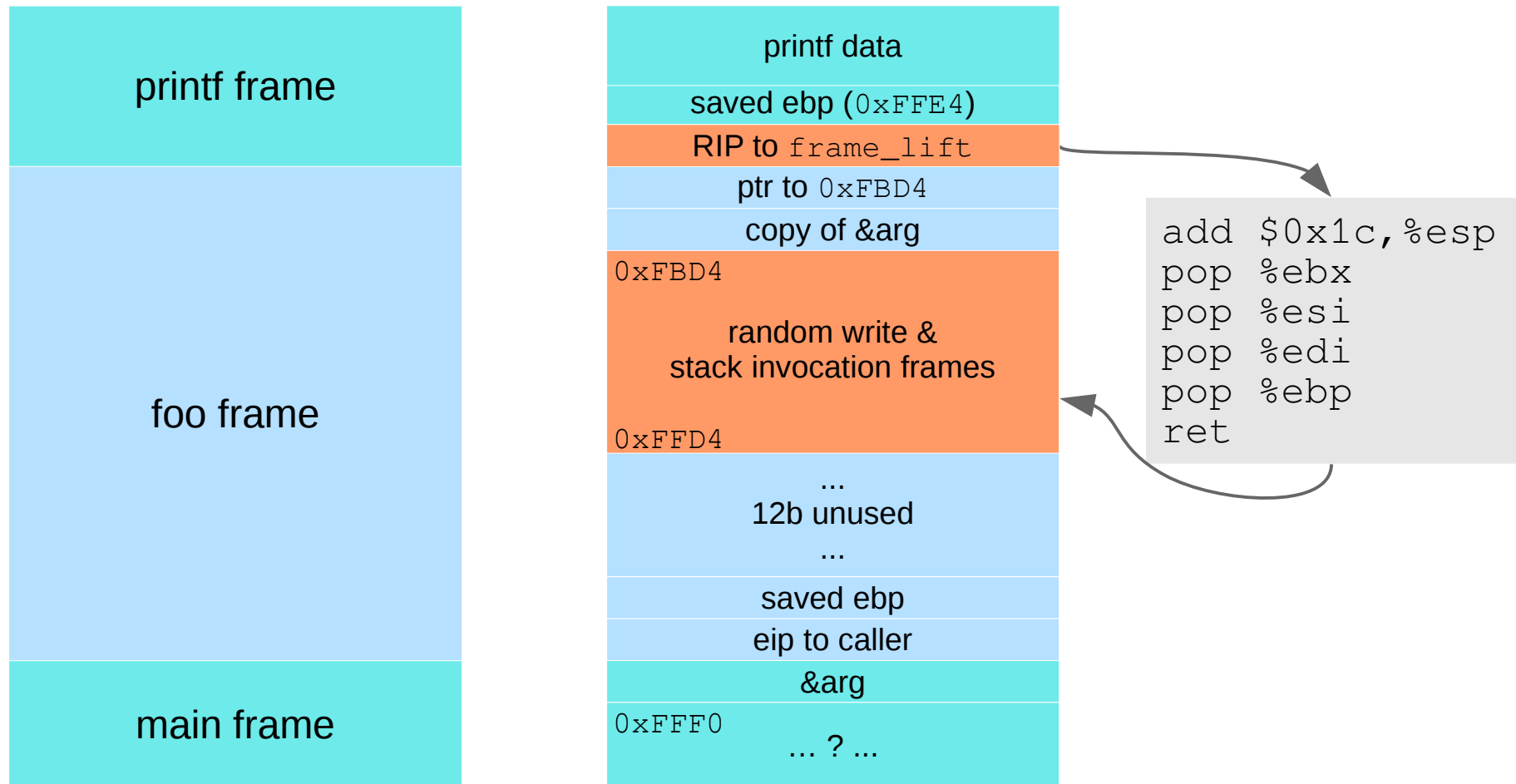

SOP: No Protection

- All addresses are known, no execution protection, no stack protection
 - Redirects control flow to code in the format string



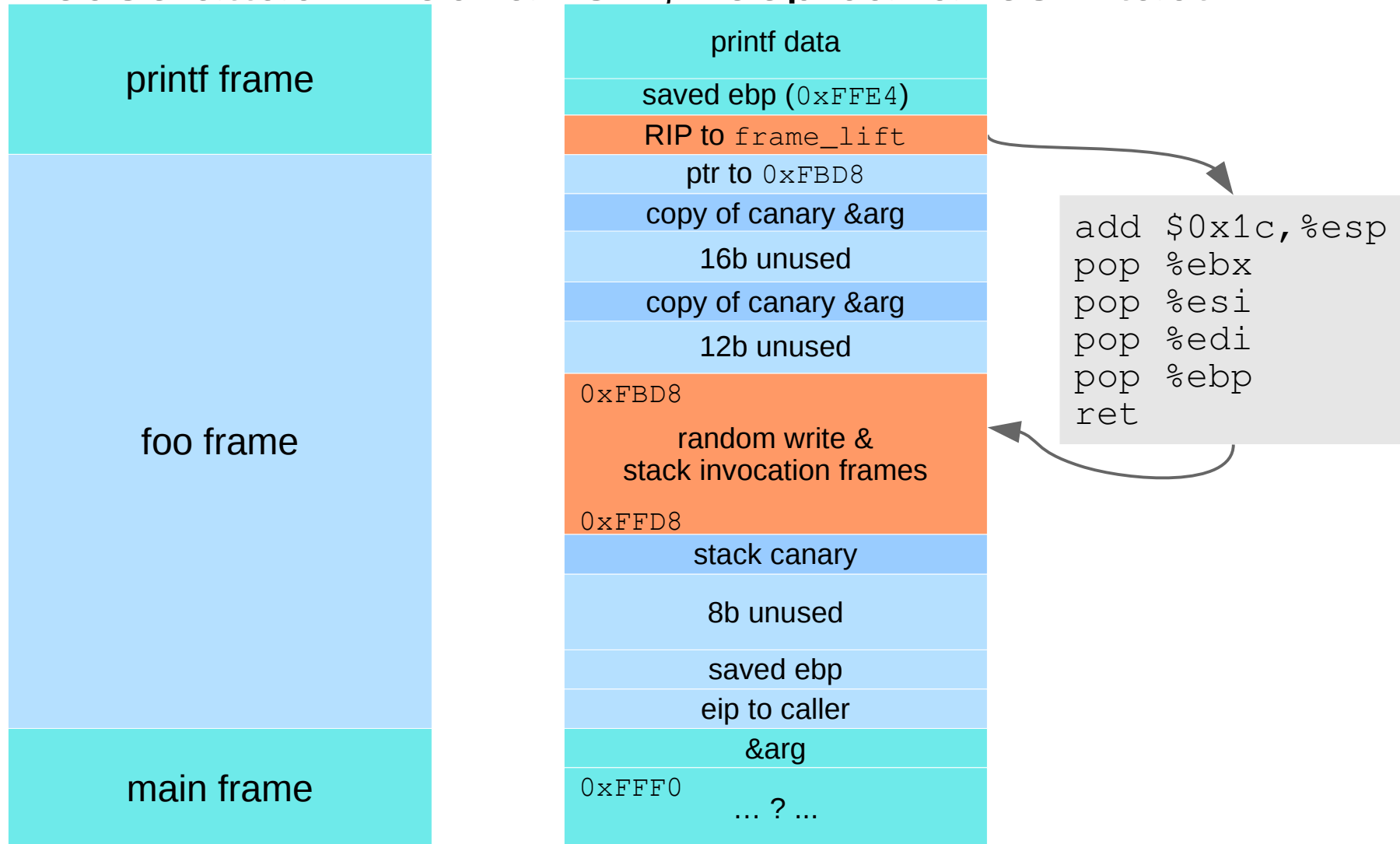
SOP: Only DEP

- DEP prevents code injection, rely on ROP/JOP instead
- GNU C compiler adds `frame_lift` gadget



SOP: DEP & ProPolice

- ProPolice uses/enforces stack canaries
 - Reuse attack mechanism, keep canaries intact



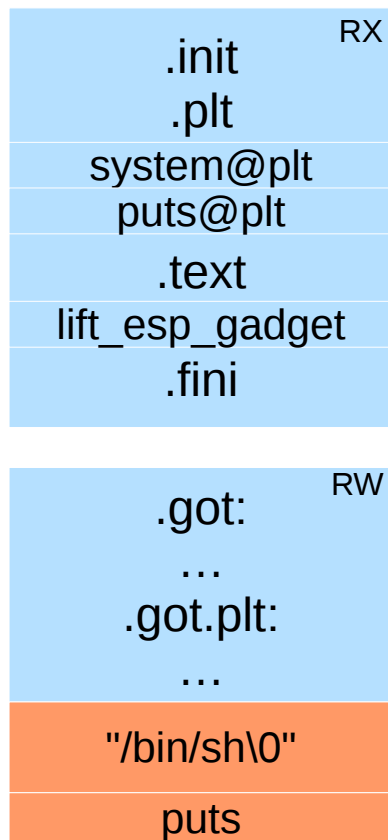
SOP: ASLR, DEP, ProPolice

- Combined defenses force SOP to reuse existing code
 - Static code sequences in the application object
 - Imported functions in the application (.plt and .got)

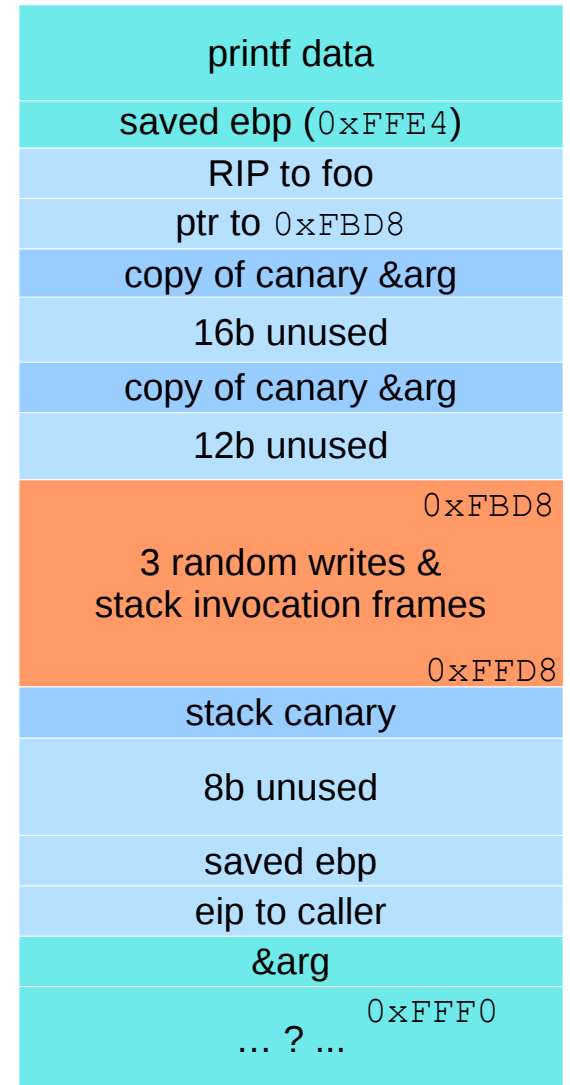
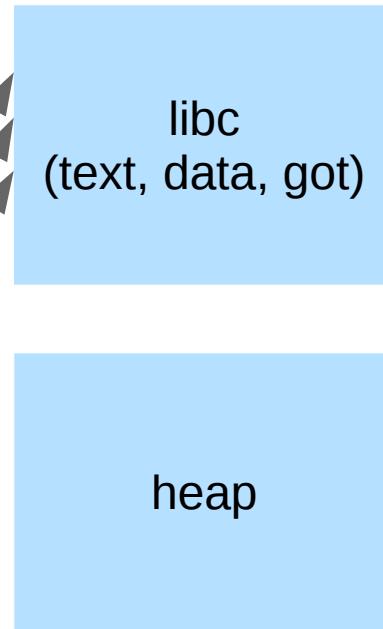
```
void foo(char *prn){
    char text[1000];           // protected on stack
    strcpy(text, prn);
    printf(text);             // vulnerable printf
    puts("logged in\n");      // 'some' function
}
```

SOP: ASLR, DEP, ProPolice

Application (static)



Libraries, heap, stack(s) (dynamic)



- Place data in RW section
- Redirect imported function (JOP)
- Use ROP for fun & profit

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Conclusion

- String Oriented Programming (SOP)
 - Relies on format string exploit
 - Extends data oriented programming (ROP / JOP)
 - Naturally circumvents DEP and ProPolice
 - Reconstructs pointers and circumvents ASLR
- Format string bugs result in complete compromise of the application and full control for the attacker
 - Protection against SOP needs more work (virtualization?)
 - Look at the complete toolchain