

# CS510 Assignment #3 (Due 3/26 before class)

March 12, 2015

## 1 Proof

(20p) Prove the following sequents:

- (a)  $p \rightarrow q, r \rightarrow s \vdash p \vee r \rightarrow q \vee s$
- (b)  $(p \wedge q) \vee (p \wedge r) \vdash p \wedge (q \vee r)$
- (c)  $\vdash \neg p \rightarrow (p \rightarrow (p \rightarrow q))$
- (d)  $q \vdash (p \wedge q) \vee (\neg p \wedge q)$  using LEM
- (e)  $p \wedge q \vdash \neg(\neg p \vee \neg q)$

## 2 Validity

(15p) Determine the validity of the following semantic entailment by constructing the truth table.

$$\neg r \rightarrow (p \vee q), r \wedge \neg q \models r \rightarrow q$$

## 3 CNF and SAT solving

(25p)

If I study, then I will not fail basket weaving 101. If I do not play cards too often, then I will study. I failed basket weaving 101. Therefore, I played cards too often.

- (a) Model the above statement to a formula. Turn the formula to its CNF form to decide validity. (15p)
- (b) Decide its validity by formulating it as a satisfiability problem and solving it. Please show the parse tree and the value assignment process. (10p)

## 4 SAT solving

(20p)

- (a) Find a formula such that the cubic SAT solving algorithm cannot decide its satisfiability, that is, it fails to converge on a deterministic value for any unmarked node. Please do not use examples from the text book. (10p)
- (b) Devise a try-and-backtrack SAT solving algorithm that operates on formulas with  $\neg$  and  $\wedge$ . You can use the linear algorithm as a primitive to construct your algorithm, e.g. *linear\_solve()* means performing the linear inference as much as possible. (10p)

## 5 The DIMACS Format

(10p) Most fast SAT solvers require the input formula in CNF. The input CNF formula is specified in the DIMACS format. Consider the following file `sample.cnf` in DIMACS format.

```
p cnf 4 5
1 0
2 -3 0
-4 -1 0
-1 -2 3 4 0
-2 4 0
```

The first line (`p cnf x y`) says that the input is a CNF formula containing  $x$  variables and  $y$  clauses. Our example has 4 variables (1, 2, 3, 4) and five clauses. The negation of a variable is denoted by putting a minus sign in front of the variable number. Each clause is described in a line terminated by a zero. Note that 0 cannot be used as a variable number. So `sample.cnf` denotes the following CNF formula:  $1 \wedge (2 \vee \neg 3) \wedge (\neg 4 \vee \neg 1) \wedge (\neg 1 \vee \neg 2 \vee 3 \vee 4) \wedge (\neg 2 \vee 4)$ .

Express the following clause set in the DIMACS format. Use the variable name  $i$  in the DIMACS format for the variable named  $x_i$  in the the above clauses. (For example, use 4 for  $x_4$ .)

$$\begin{array}{ll} \omega_1 = x_1 \vee x_3 \vee x_4 & \omega_3 = \neg x_1 \vee x_2 \vee x_3 \\ \omega_2 = x_1 \vee \neg x_2 \vee x_3 & \omega_4 = x_1 \vee \neg x_2 \vee x_4 \\ \omega_5 = x_2 \vee x_4 & \omega_6 = x_3 \vee x_4 \end{array}$$

## 6 Using a SAT solver

(10p)

Some publicly available fast SAT solvers are `MiniSat`, `zChaff`, `siege`. For this assignment you will install and use the `MiniSat` SAT solver which was the fastest SAT solver in the SAT-competitions of 2005 and 2006. You can run `MiniSat` SAT solver simply by the following command:

```
/path/MiniSat.v1.14.linux sample.cnf sample.result
```

The file `sample.cnf` is a description of a CNF formula in DIMACS format. `MiniSat` reports whether the given formula is (un)satisfiable in the file `sample.result`. If the formula is satisfiable, then a satisfying assignment is also written to `sample.result`.

Run `MiniSat` on the DIMACS file you made from the previous problem. What is the output stored in results file?